



A Trustworthy Middlebox-aware Networking Architecture

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Motivation

To be practical, TLS with middleboxes must address not only **secure participation of middleboxes** but also **trustworthiness of middleboxes**

1 Middleboxes are valuable but become useless with TLS

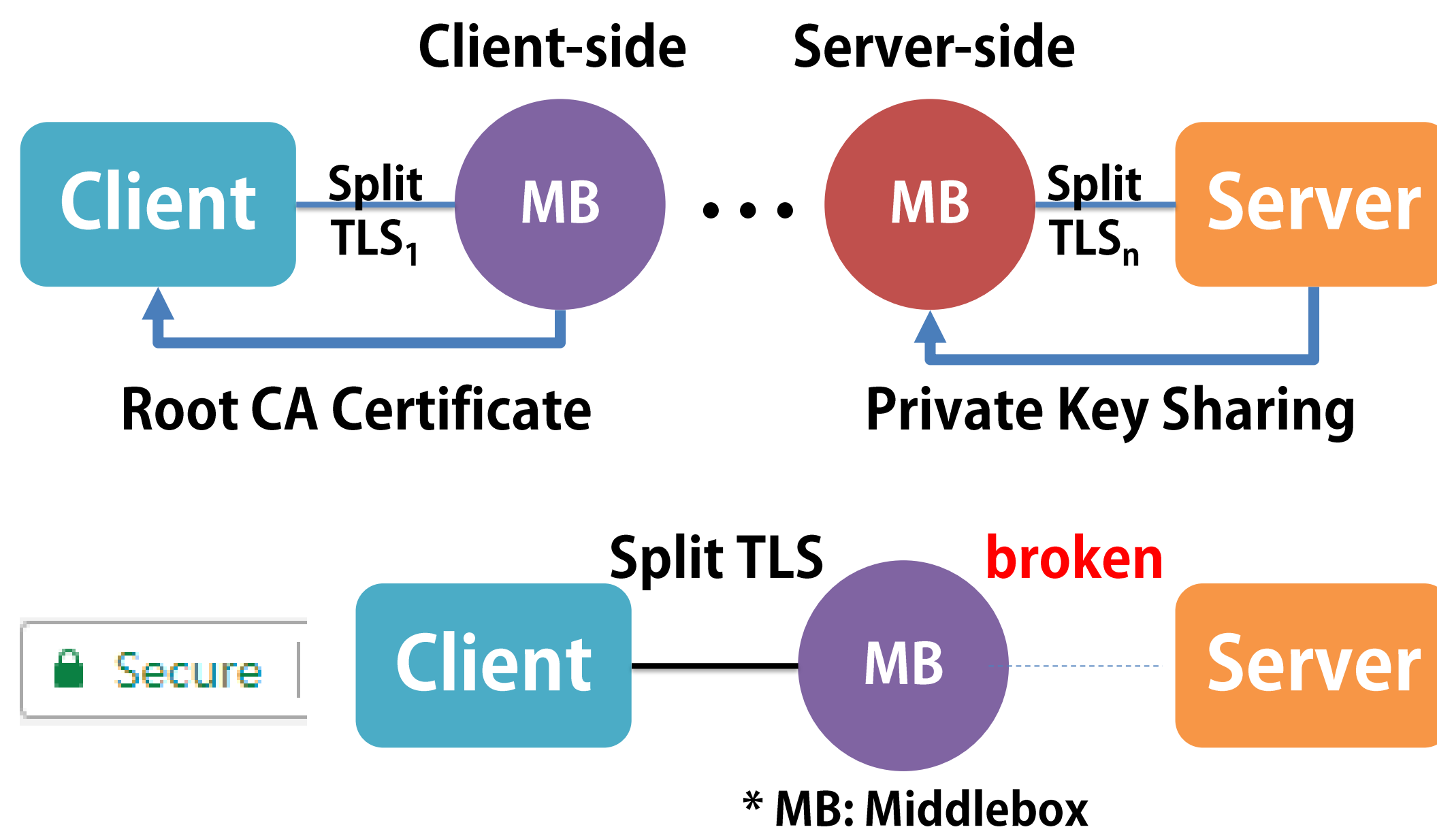


Middleboxes can offer

- Defense from attacks at the network edge
- Performance improvement

However, value-added service for security and performance cannot be carried out with TLS

2 Current TLS Interception (Split TLS) breaks security and infringes privacy



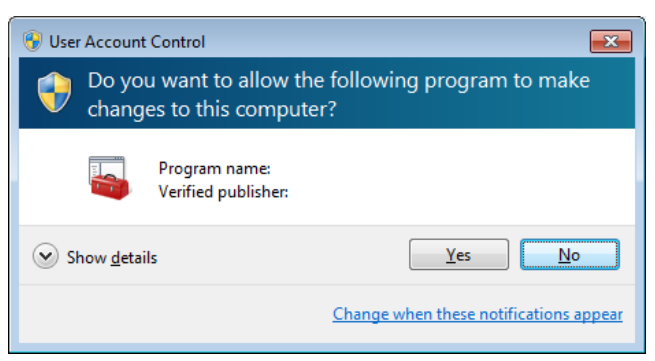
Even though data secrecy, data authentication, and entity authentication are broken, a user agent still say "secure"

3 We do not know how to trust TLS middleboxes

Much research assumes trustworthiness of middleboxes through their certificates

Only certificate validation for trust is insufficient since a user does not explicitly intend to communicate with middleboxes

User access control can be used, but it will decrease usability



Trustworthiness Building Blocks

We define trustworthiness of middleboxes as **public auditability** which becomes feasible with **explicit authentication** and **modification check**

1 Difference from a man-in-the-middle attacker

We add **auditability** on **read/write** operation to be different from passive/active adversary

For read operation,

Explicit authentication by Middlebox Certificate in Middlebox Transparency

For write operation,

Modification check with Modification Record

2 Middlebox Certificate in Middlebox Transparency

Middlebox Certificate includes

- Name of a middlebox provider
- Read/Write permissions on key usage
- Role of a middlebox
- Access URL to information page of a middlebox

All CAs or middlebox providers must register middlebox certificate into middlebox transparency, the same **public log server** with certificate transparency

Only middlebox having middlebox certificate can participate into TLS session as a **reader** and every entity can **audit** readers in middlebox transparency

3 Modification Record

This is a data structure recording the **history of modification** by middleboxes, starting from a source

Writer records its ID, the hash of original message, and HMAC of modification whenever it modifies

$$m' \quad H(\text{secret}_s, H(m)) \quad ID_{mb} \quad H(m) \quad H(\text{secret}_{mb}, H(m) || H(m'))$$

Example: Server sends m and MB modifies it into m'

Experiment shows verification only takes less than 350 us on the desktop, even with 100 writers

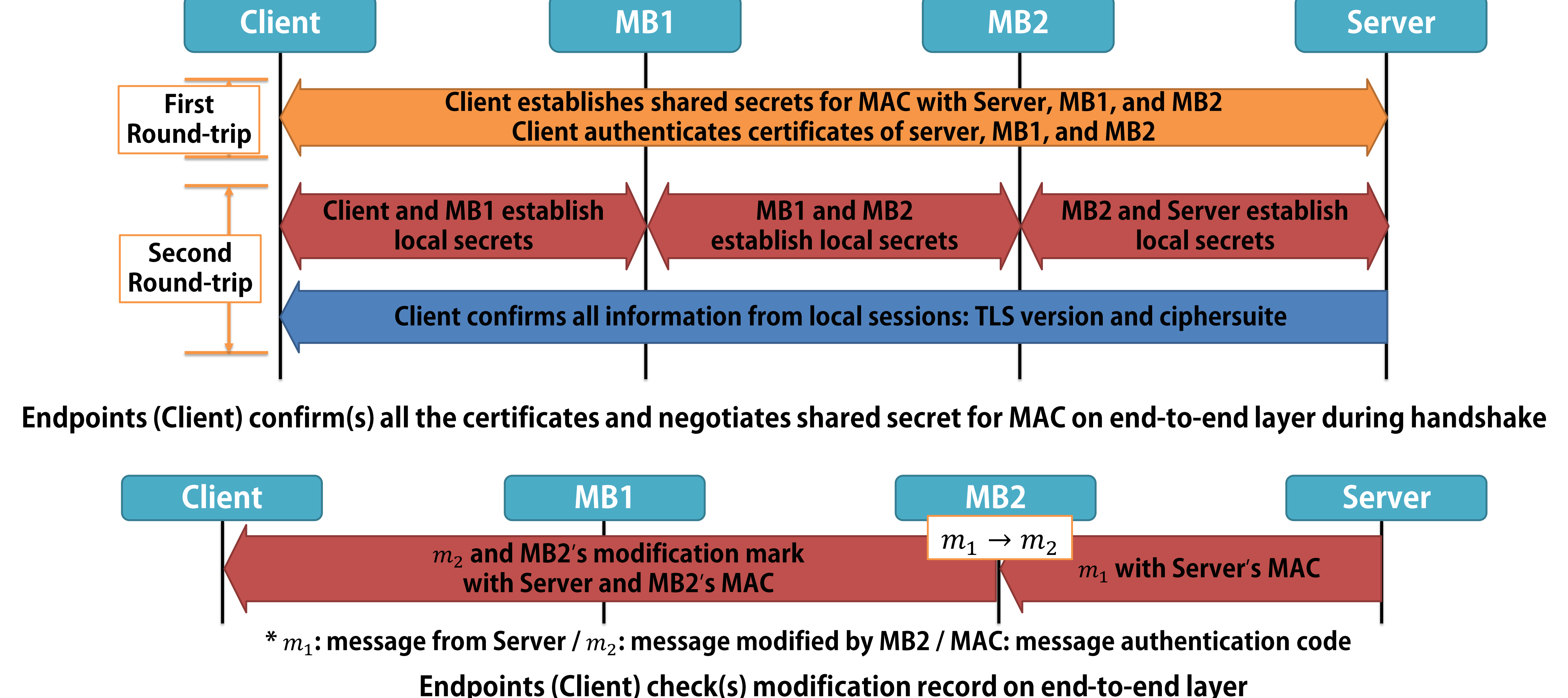
Middlebox-aware TLS

Middlebox-aware TLS guarantees **new TLS security requirements** with **trustworthiness-managed middleboxes**

1 New Security Requirements

- Data secrecy → **Path secrecy**
Endpoints (Client) confirm(s) every local session is encrypted with higher TLS version/ciphersuite
- Data authentication → **Data source authentication**
Endpoints (Client) confirm(s) the source of data, that is, who sends the data
→ **Modification accountability**
Endpoints (Client) confirm(s) valid modification, that is, who modifies the data
- Entity authentication → **Server/Middlebox authentication**
Endpoints (Client) authenticate(s) all the entities in the session by their certificates

2 End-to-end TLS layer over Split TLS layer



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